

Computer Architecture

Dr. Esti Stein

(Partly taken from Dr. Alon Schclar slides)

Based on slides by:

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Department of Information & Communication

Taken from: **M.**

**Mano/Computer Design and
Architecture 3rd Ed.**

Course Administration

- **Instructor**

Carmi Merimovich (carmi@mta.ac.il)

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- **Exercise checker**

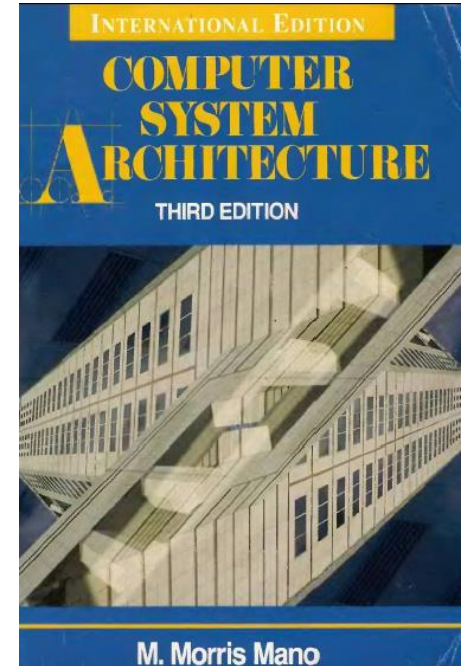
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- **Books**

Mano M.: **Digital Design**, Prentice Hall.

Patterson and Hennessy, [Computer Organization Design, The Hardware/Software Interface](#), 3rd Edition, Morgan Kaufmann.



Course Administration

- Grade
 - Exam – 100% (must pass > 60)
 - Exercises
 - Every 1-2 weeks
 - 100% mandatory submission – 10 point reduction for each overdue day
- Office:
 - Zoom meetings
- Office hour
 - Email us

Course Outline

פרק 4	מבוא כללי ושפת RTL	שבוע 1
פרק 5	ארגון המחשב – ה-BUS	שבוע 2
פרק 5	ארגון המחשב – יחידת הבקרה	שבוע 3
פרק 5	ארגון המחשב – פסיקות	שבוע 4
פרק 6	תכנות המחשב הבסיסי	שבוע 5
פרק 6	תכנות המחשב הבסיסי	שבוע 6
פרק 6	תכנות המחשב הבסיסי	שבוע 7
פרק 7	מיקרו תכנות	שבוע 8
פרק 7	מיקרו תכנות	שבוע 9
פרק 11	יחידות הקלט/פלט	שבוע 10
פרק 11	יחידות הקלט/פלט	שבוע 11
פרק 12	מבנה הזיכרון	שבוע 12
פרק 12	זיכרון CACHE	שבוע 13

What is “Computer Architecture”

Architecture:

Design a building that
is well suited to its
purpose



Computer
Architecture:

Design a computer
that is well suited
to its purpose



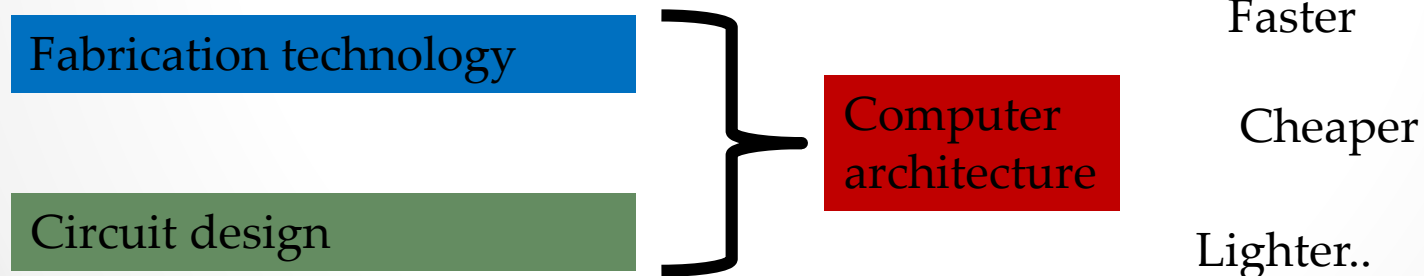
Why do we need computer architecture

Improve performance:

speed, battery life, size, weight, energy efficiency..

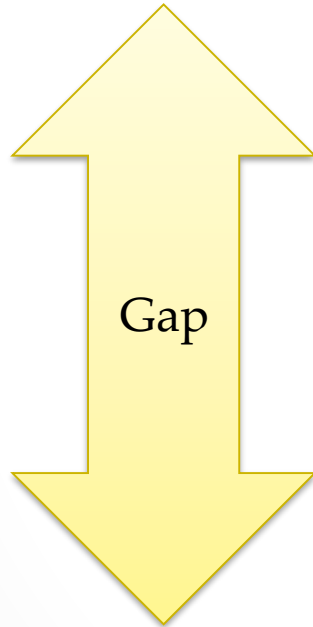
Improve abilities:

3d graphics, debugging support, security..



From Application to Physics

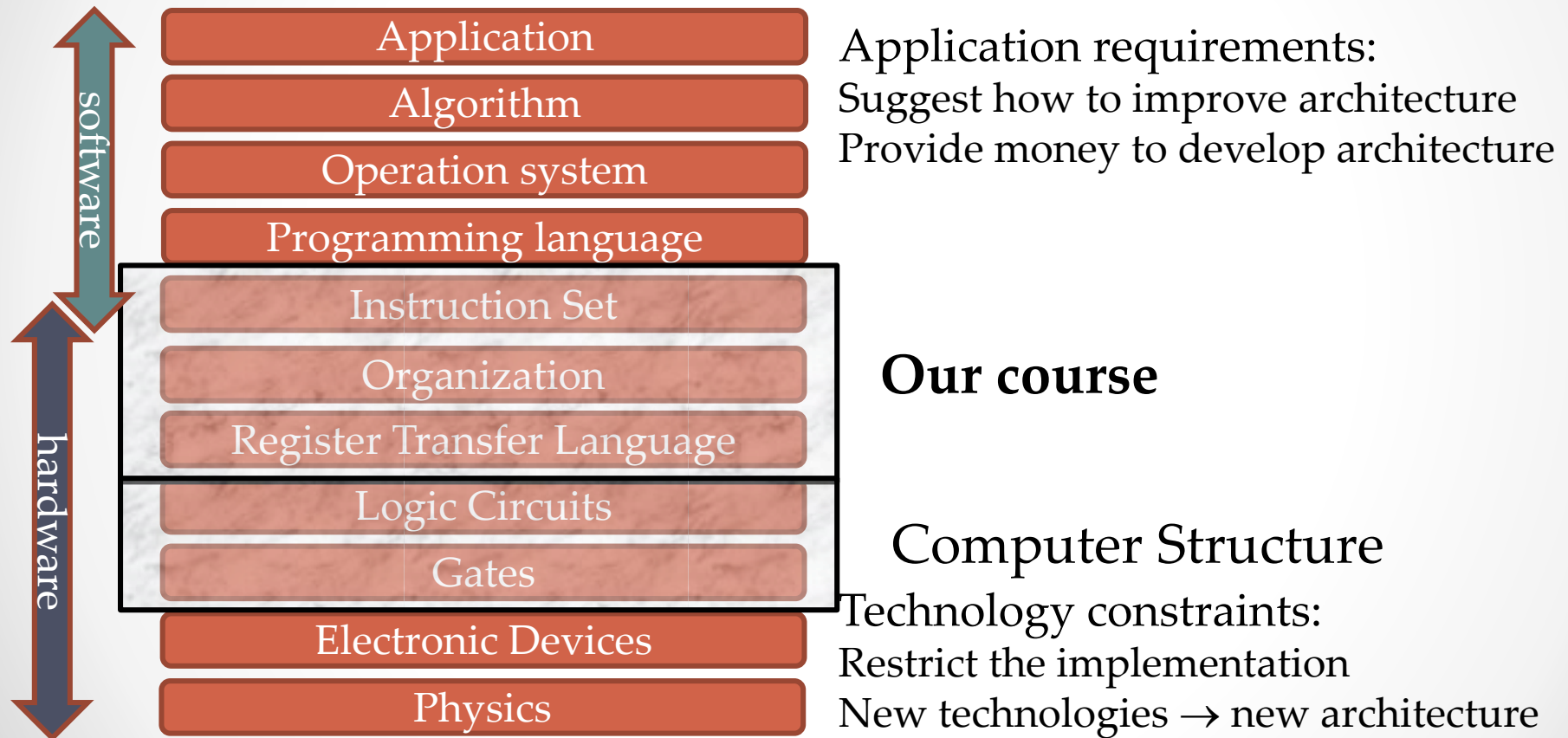
Application



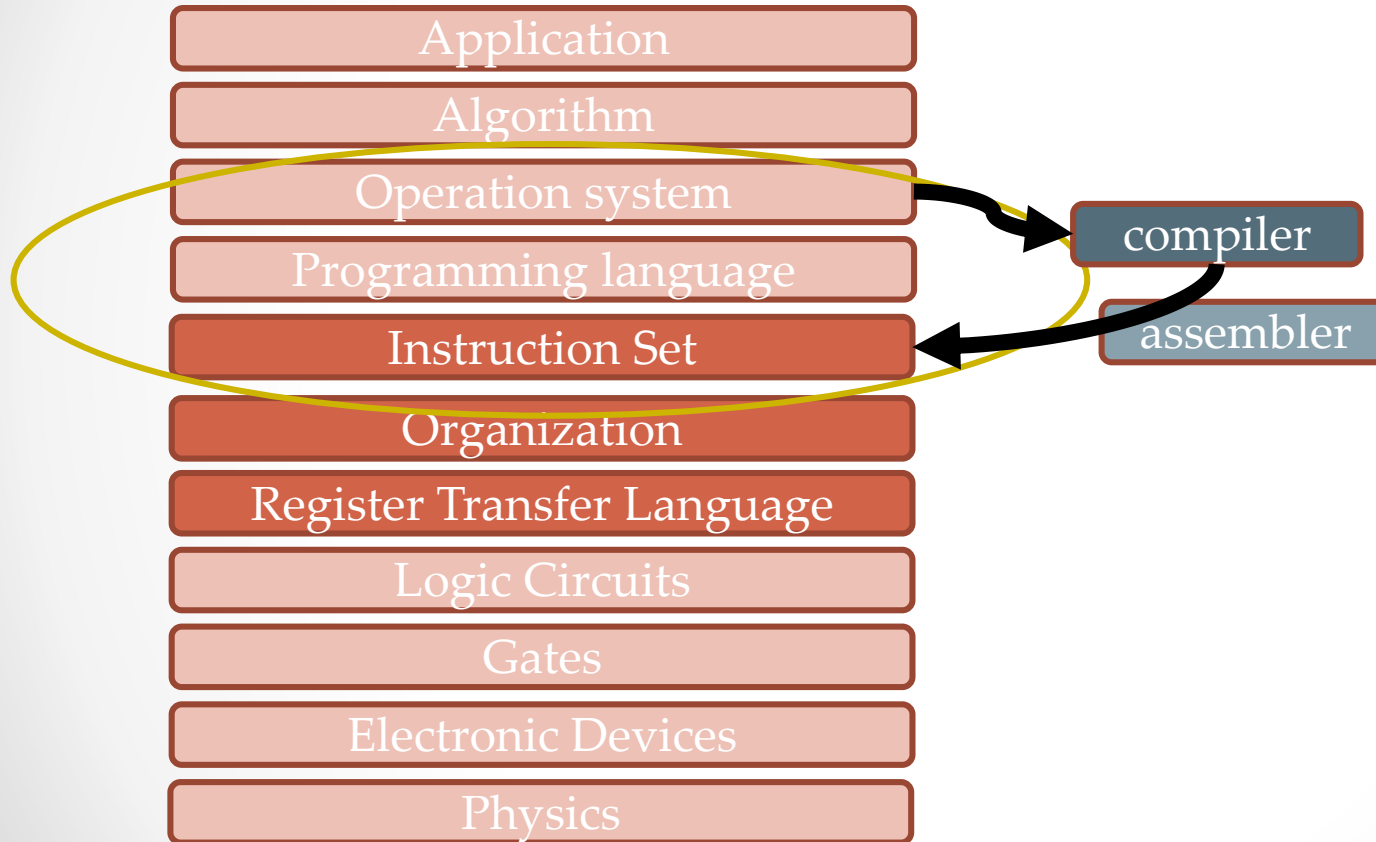
Physics

It's the abstraction of the implementation layers that needs to bridge the gap

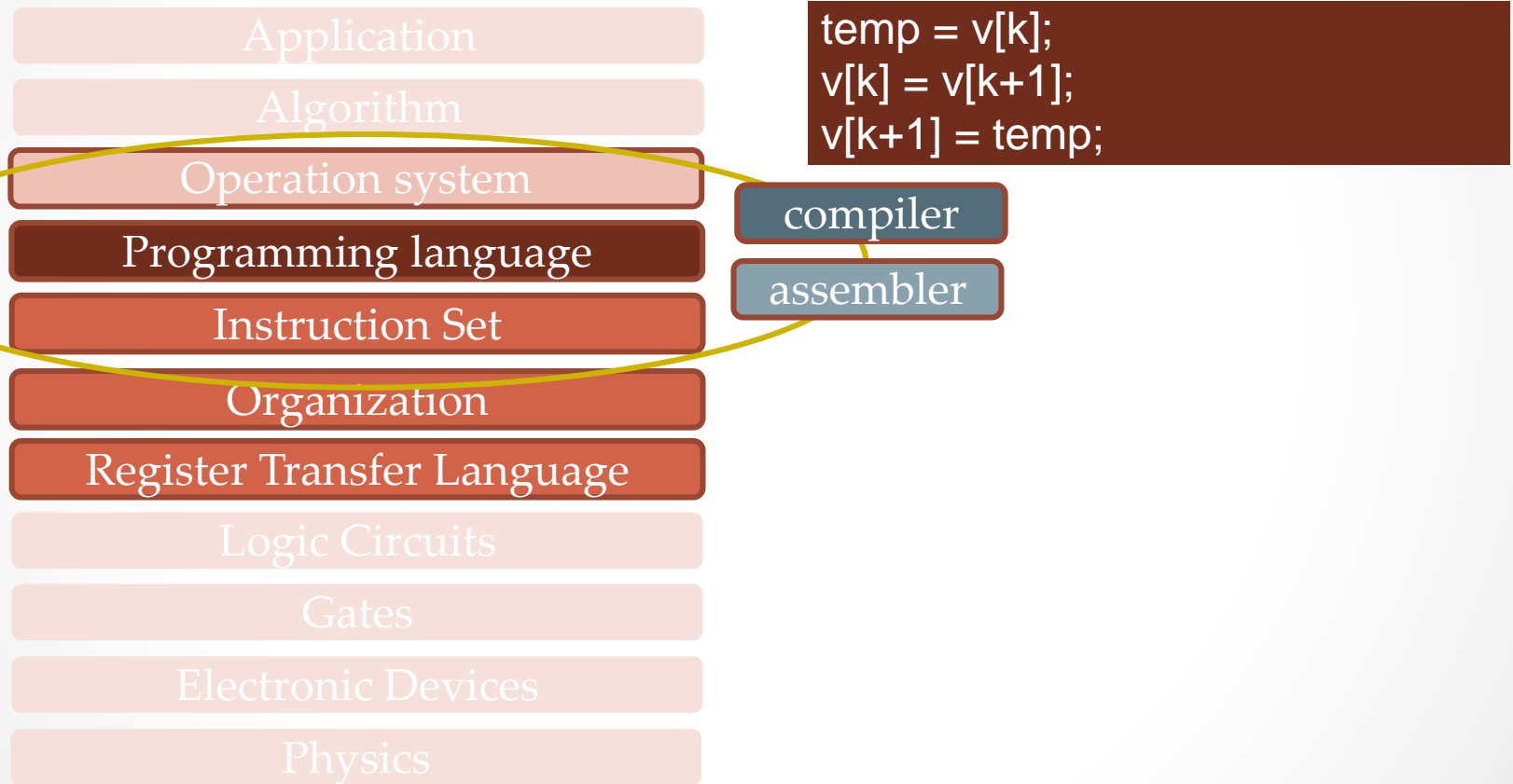
From Application to Physics



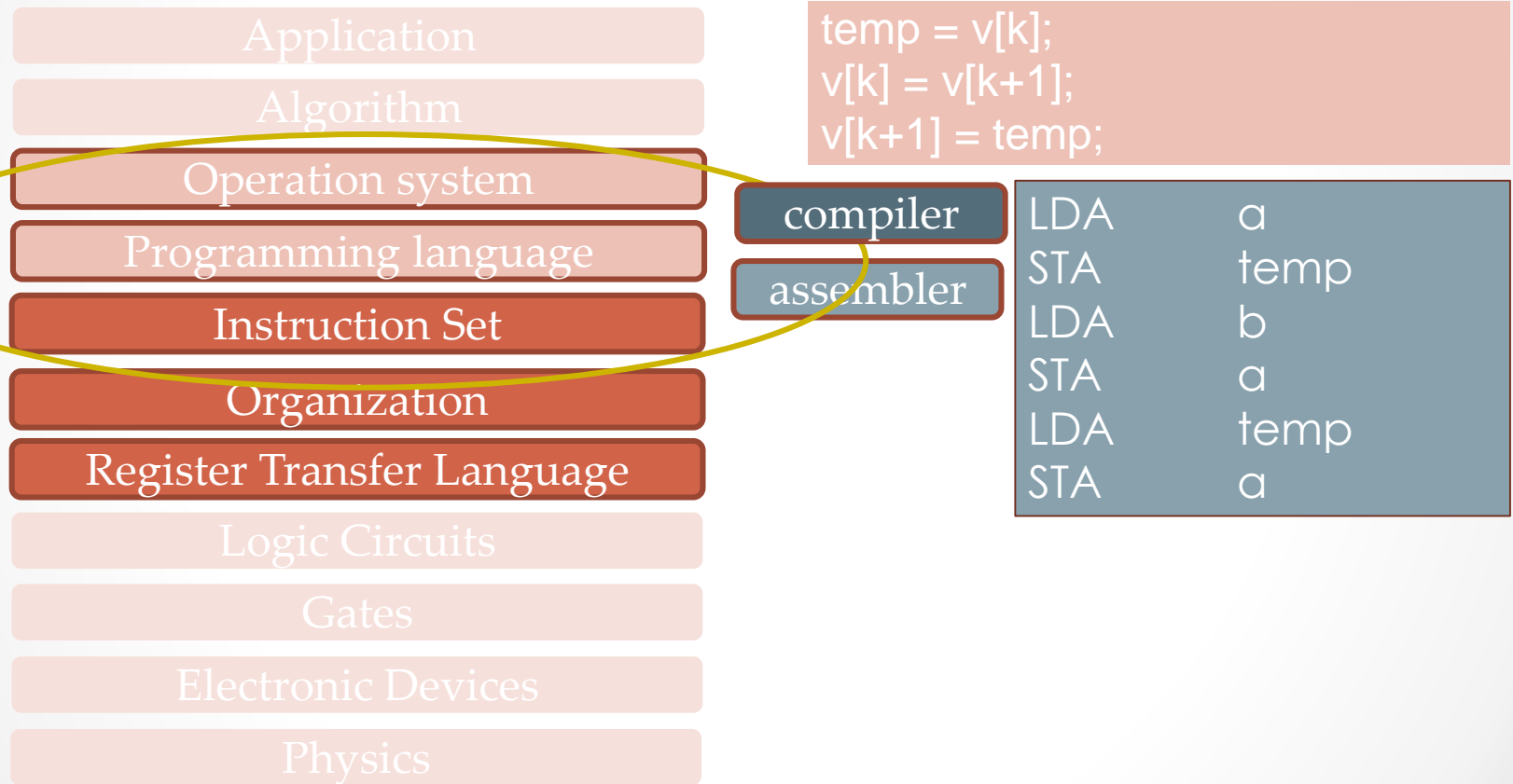
From Application to Physics



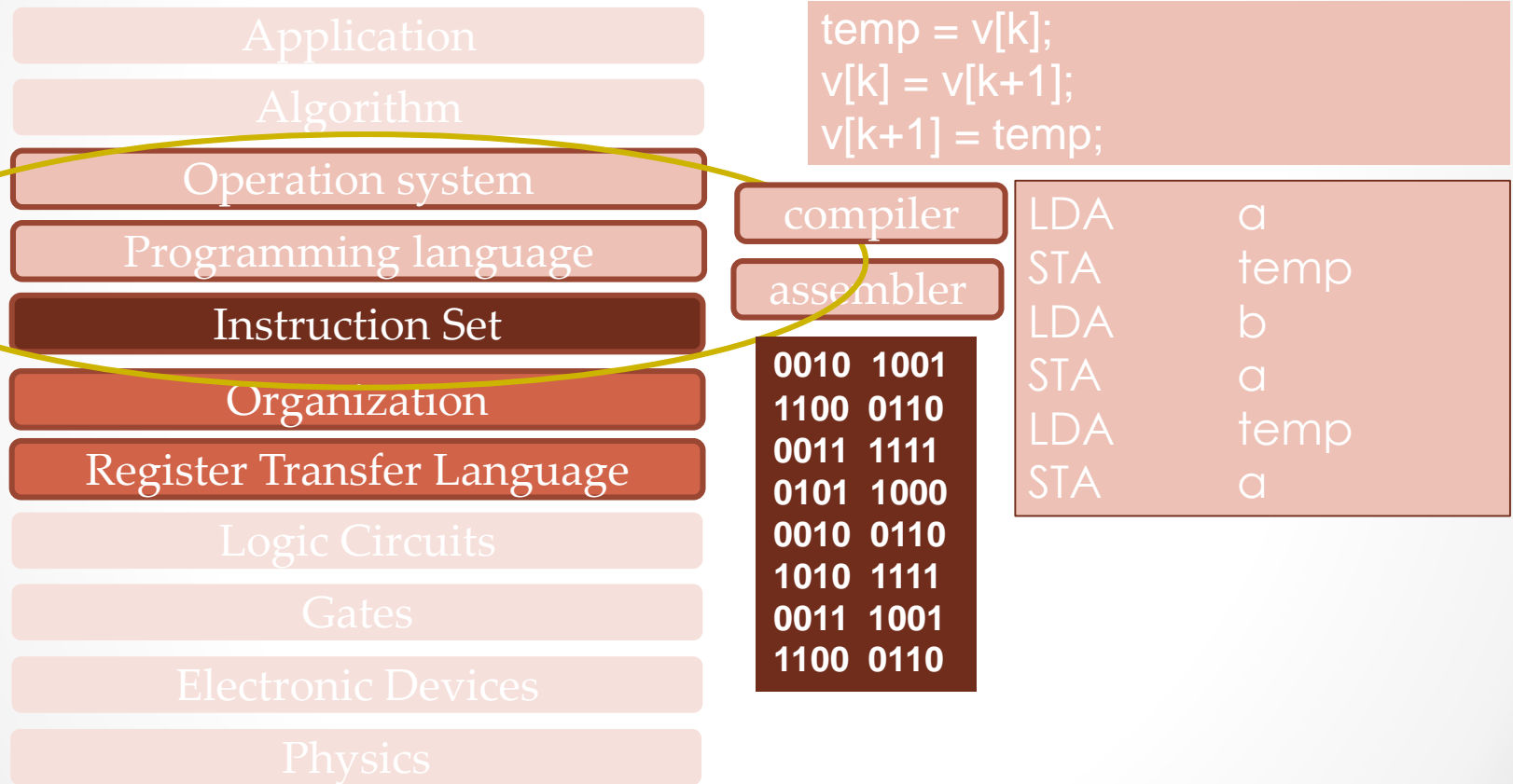
From Application to Physics



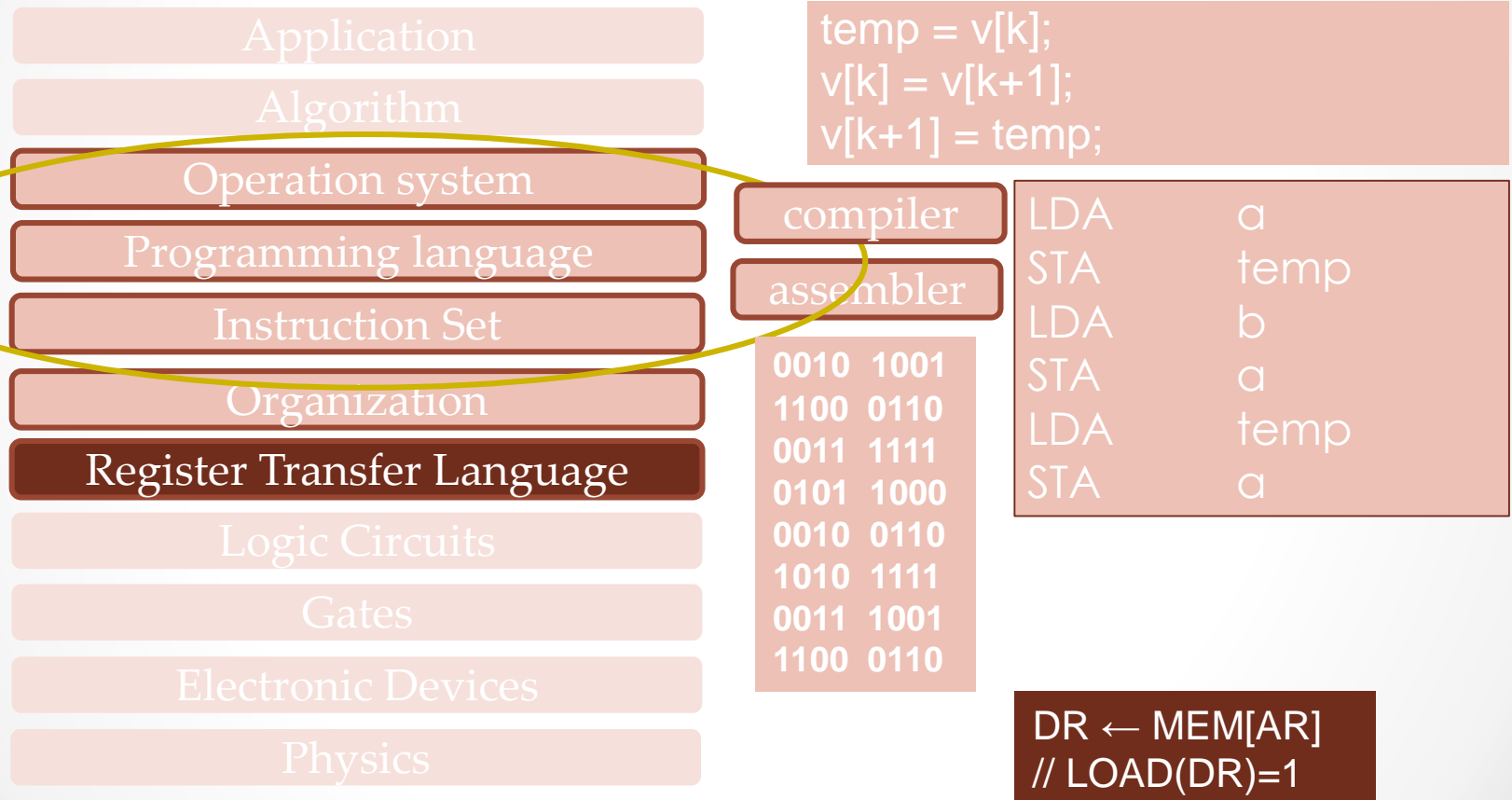
From Application to Physics



From Application to Physics



From Application to Physics



Talking to the Computer

- In order to “**talk**” with the computer we must send it electronic signals.
- The easiest signals for an electronic machine to understand are ***on*** and ***off***
 - Correspond to *high* voltage and *low* **voltage**.
- Thus the computer’s alphabet is composed of two symbols **0** and **1**.
- Any “words” composed of these 2 numbers are called ***binary numbers***.

Talking to the Computer

- A computer needs **our instructions** in order to function
- Instructions are formulated as **binary** numbers
- The binary number **11101100100001** can be an instruction to **subtract** two numbers
- Every computer understands a predefined set of instructions – **instruction set**
- An **instruction** has a **predefined structure** which matches the architecture of the computer

Assembly Language

- Initially (in the late 40s) computer **programs** were written as **binary numbers**
- They were input to the computer by **turning on and off** switches

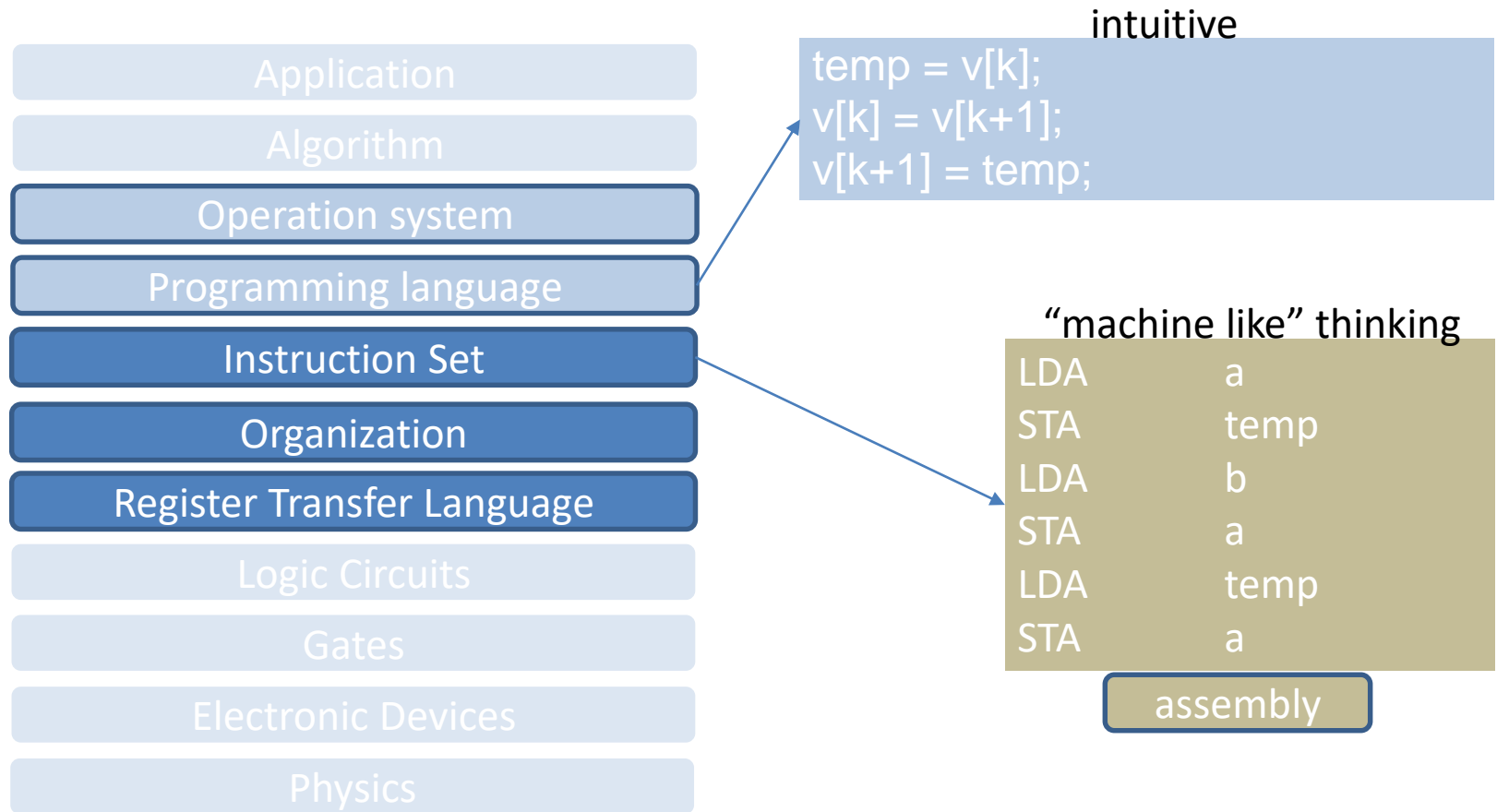


- It becomes extremely **tedious** after a short while

Assembly Language

- Next step: **automatic conversion** of the codes into binary instructions
- This program is called an **assembler**.
- The symbolic names of the instructions are called the **assembly language** :
 - Increases productivity
 - However
 - Each machine instruction must be written in a single line
 - The programmer needs to think like a machine
- A higher level of abstraction is needed

High-level & Assembly Languages



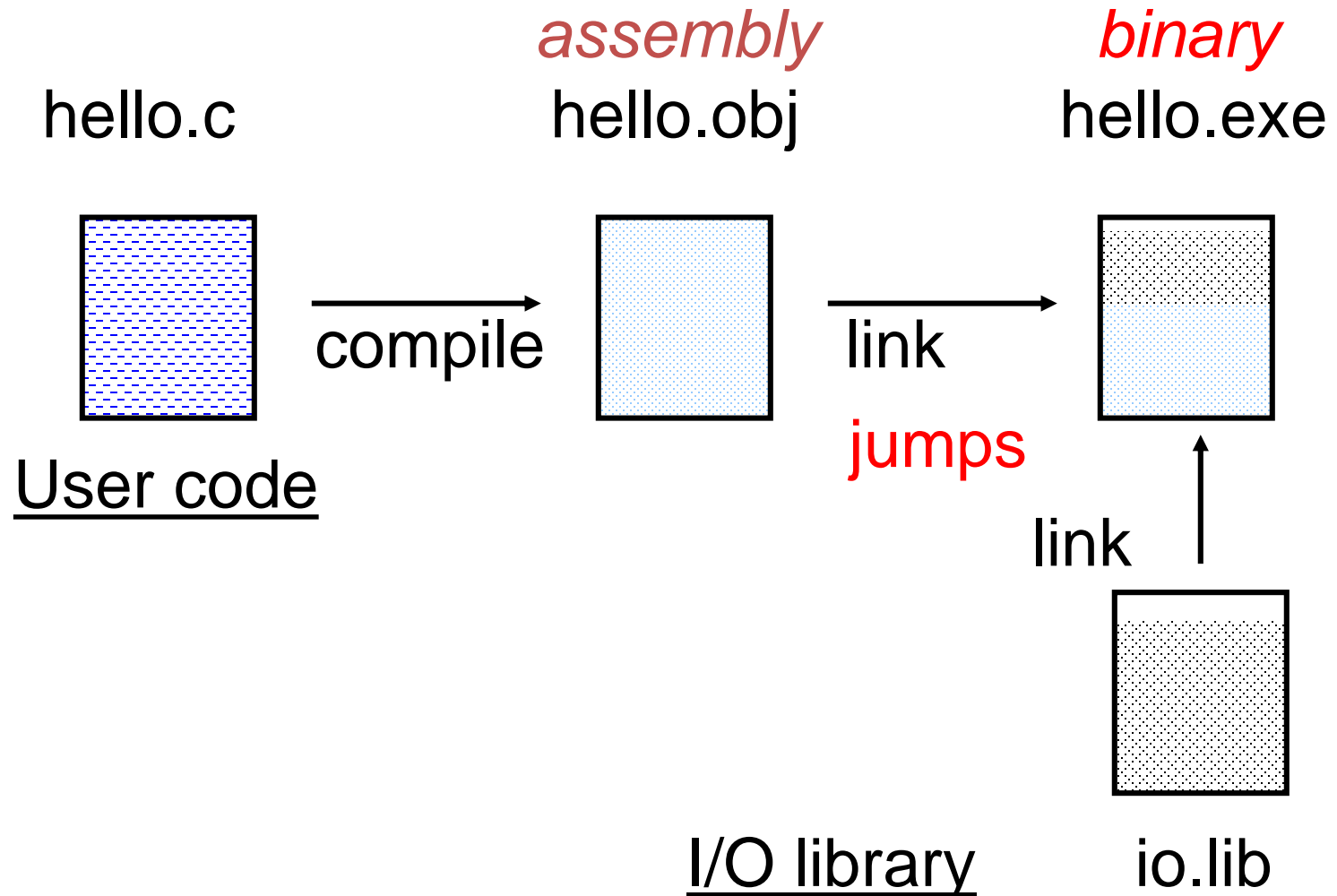
High Level Languages

- Humans think in a language that consists of **sentences** over the alphabet
- Why not **define** a **language** that is more **intuitive** for the programmer
- Next step :
 - Define a **high level language** (Fortran, JAVA, C)
 - Construct a translator that converts high level instructions into binary instructions – the **compiler**

A Simple Example

- The high level expression in C:
A = A + B
- Translates into the assembly instruction
LDA A
ADD B
STA A
- Translates into the binary instruction:
0010 110010100000
0001 100010100001
0011 110010100000

Compiling and Linking a Program



Executable program (EXE)

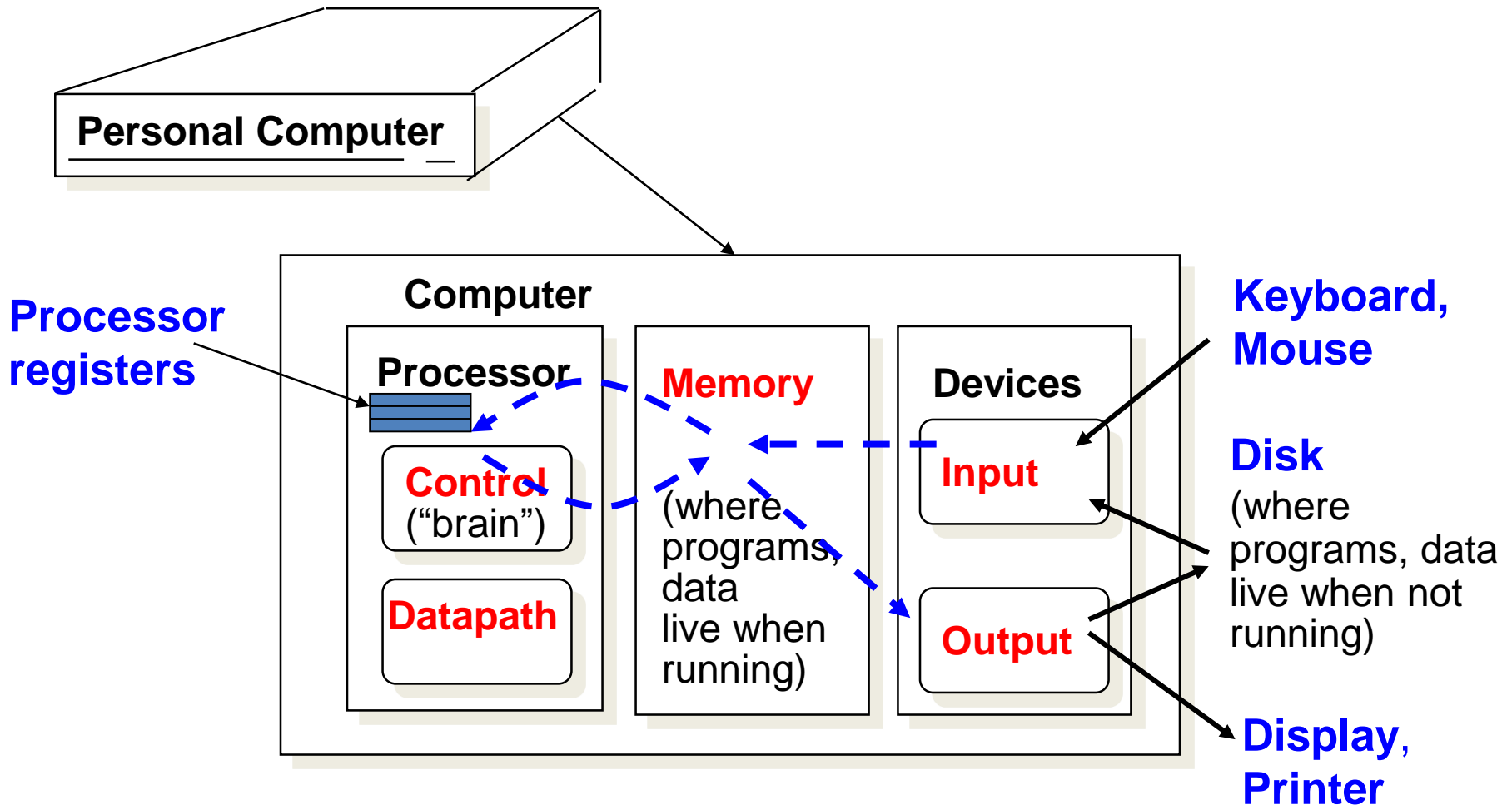
- A sequence of machine instructions
 - Binary coded
 - Operations, operands (values)
 - Operands may be
 - Memory addresses (= variables)
 - Values – floating point, 2's-comp
 - Register numbers

From EXE file to execution

- Given a binary executable file
 - the file is ***loaded*** from the I/O device (hard disk, DVD, Disk-on-Key etc.) into the memory
 - the processor **executes** the instructions in an **iterative** process
 - the **operating system** coordinates this process

Anatomy:

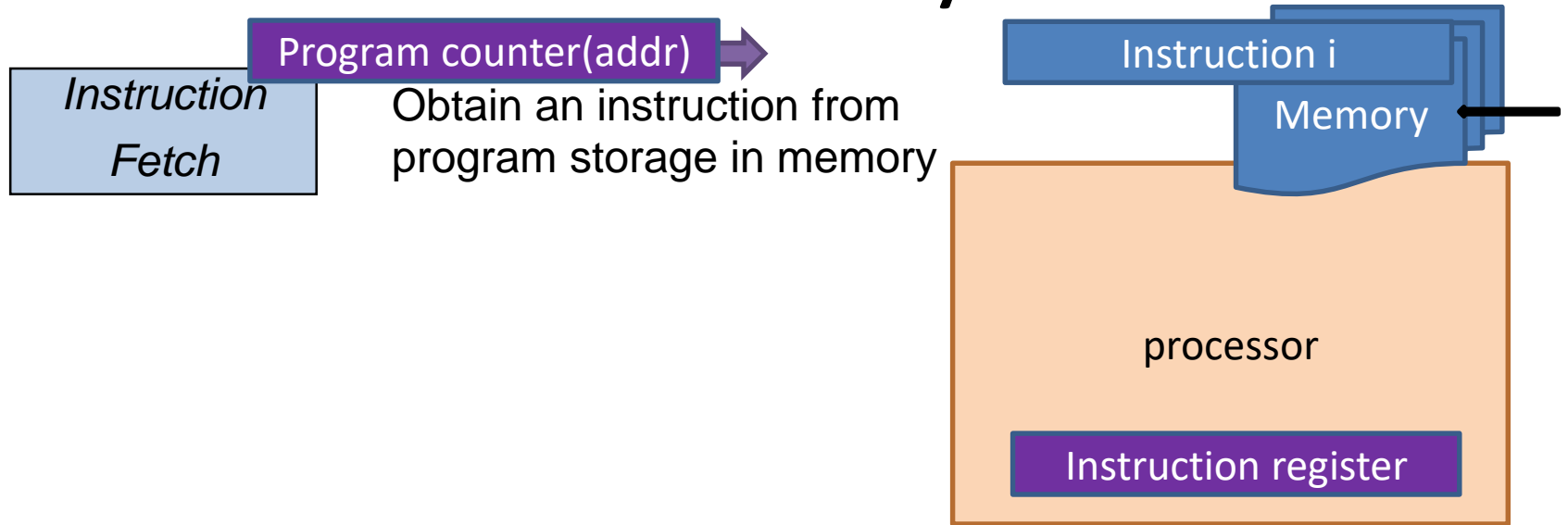
5 components of any Computer



Inter-component flow

- Data flows from the input devices into memory
- From the memory into the processor
- The data is processed and written back to memory
- It is then stored or displayed in the output devices

Execution Cycle



Execution Cycle

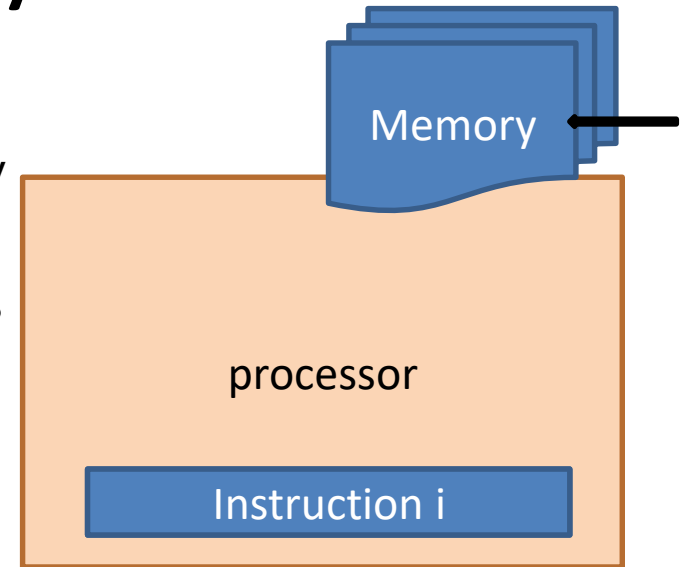
*Instruction
Fetch*



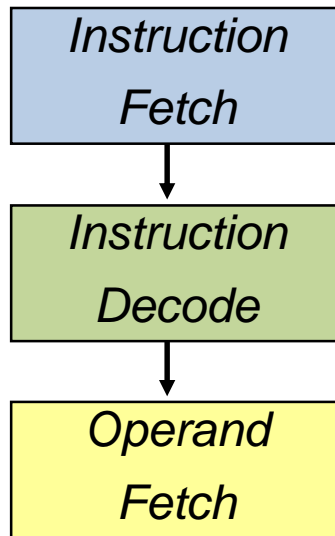
*Instruction
Decode*

Obtain instruction from
program storage in memory

Determine required actions
and instruction size



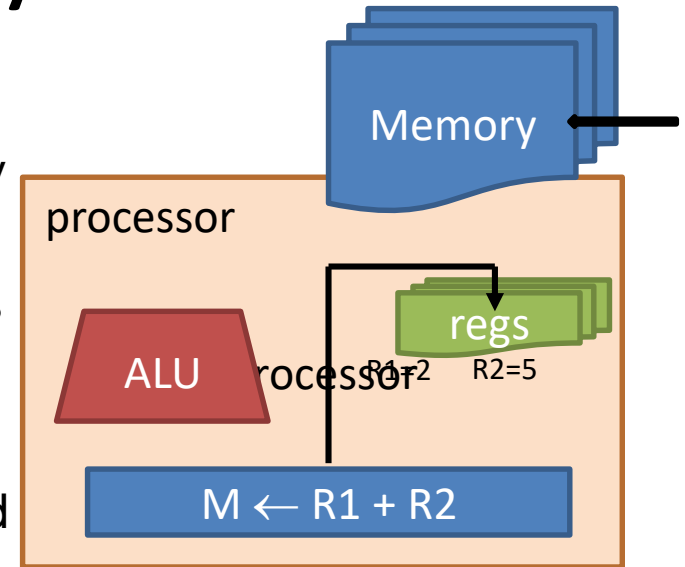
Execution Cycle



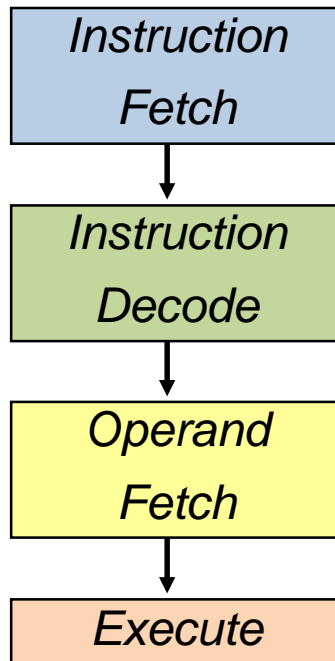
Obtain instruction from program storage in memory

Determine required actions and instruction size

Locate and obtain operand data



Execution Cycle

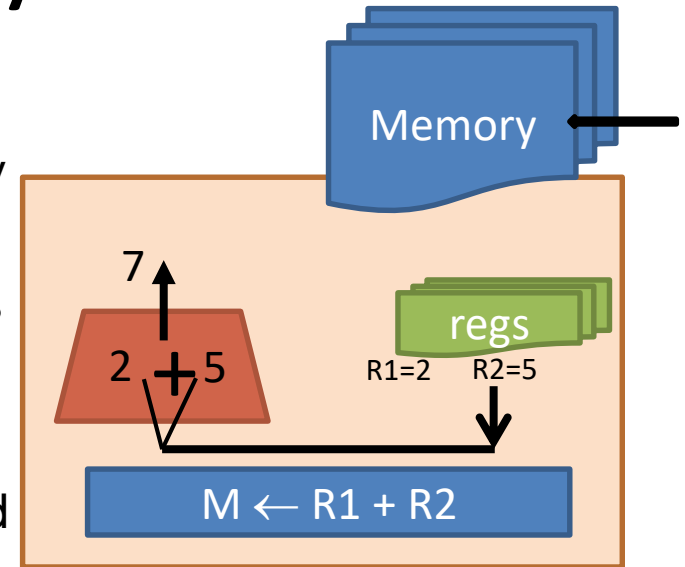


Obtain instruction from program storage in memory

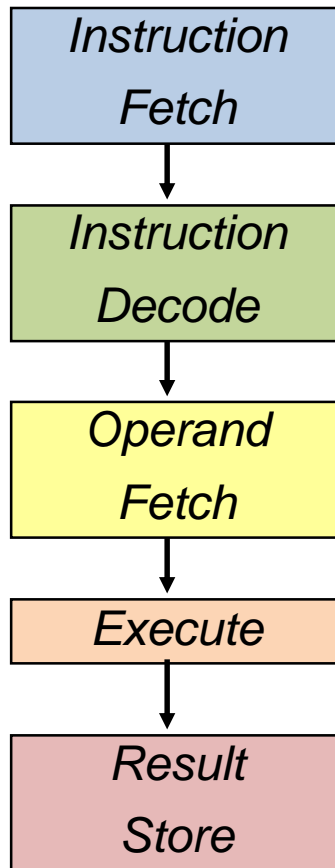
Determine required actions and instruction size

Locate and obtain operand data

Compute result value or status



Execution Cycle



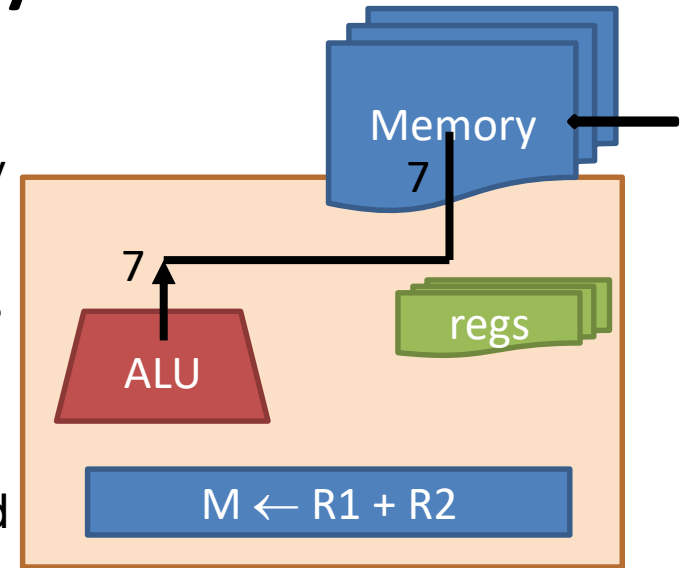
Obtain instruction from program storage in memory

Determine required actions and instruction size

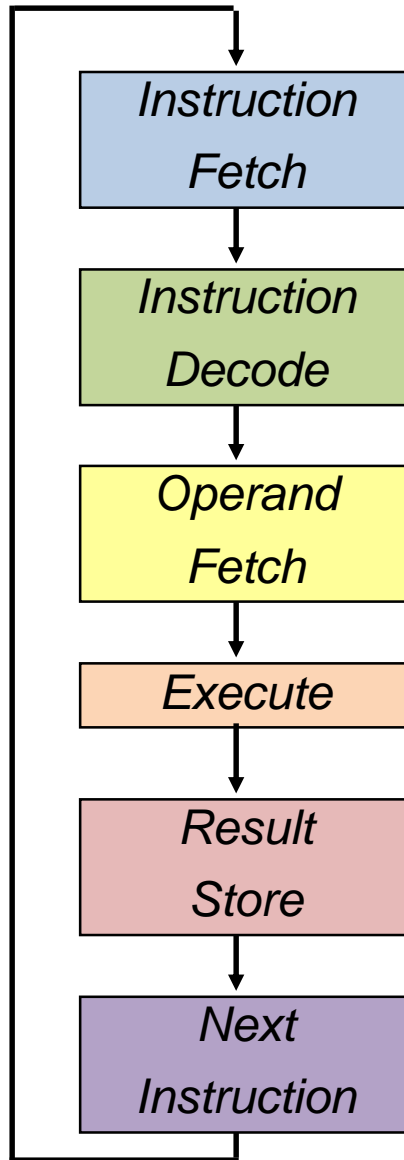
Locate and obtain operand data

Compute result value or status

Deposit results in storage



Execution Cycle



Obtain instruction from program storage in memory

Determine required actions and instruction size

Locate and obtain operand data

Compute result value or status

Deposit results in storage

Determine next instruction (not the next in case of **branch**)

